

**Performanse računarskih sistema (IR4PRS, SI4PRS)  
-rešenja zadataka-**

1. Videti predavanja.

2.

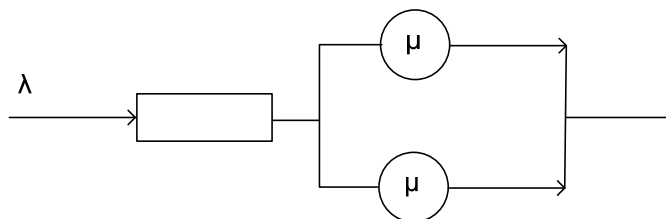
$$\begin{aligned} \overline{T_{am}} &= \frac{2}{600^2} \int_0^{600} (600-x) \cdot T_{am}(x) \cdot dx = \\ &= \frac{2}{600^2} \cdot \left( \int_0^{400} (600-x) \cdot 0.5 \cdot x \cdot dx + \int_{400}^{600} (600-x) \cdot 10\sqrt{x} \cdot dx \right) = \\ &= \frac{2}{600^2} \cdot \left( \int_0^{400} \left(300x - \frac{x^2}{2}\right) \cdot dx + \int_{400}^{600} (6000\sqrt{x} - 10x^{3/2}) \cdot dx \right) = \\ &= \frac{2}{600^2} \left( 300 \cdot \frac{1}{2} x^2 \Big|_0^{400} - \frac{1}{6} x^3 \Big|_0^{400} + 6000 \cdot \frac{2}{3} x^{3/2} \Big|_{400}^{600} - \frac{20}{5} x^{5/2} \Big|_{400}^{600} \right) = \\ &= \frac{2}{600^2} \left( 150 \cdot 400^2 - \frac{1}{6} 400^3 + 4000 \cdot (600^{3/2} - 400^{3/2}) - \frac{20}{5} (600^{5/2} - 400^{5/2}) \right) = \\ &= \frac{1}{18} \left( 150 \cdot 16 - \frac{1}{6} 400 \cdot 16 + 400 \cdot (6^{3/2} - 4^{3/2}) - 40 \cdot (6^{5/2} - 4^{5/2}) \right) = \\ &= 98.047ms \end{aligned}$$

$$\overline{T_{rd}} = \frac{1}{2} T_{rev} = \frac{30}{N_{rev}} = 5.555ms$$

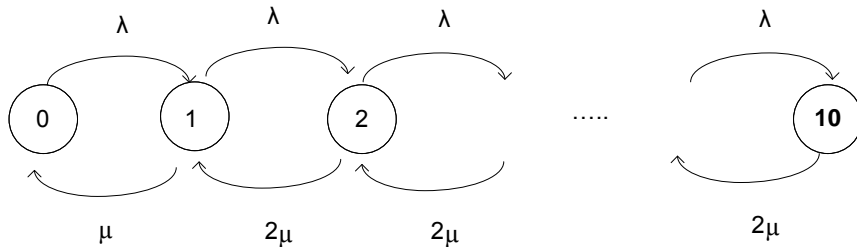
$$T_{dt} = \frac{1}{8} T_{rev} = \frac{5}{N_{rev}} = 1.389ms$$

$$\overline{T_{uk}} = 1200 \cdot (\overline{T_{am}} + \overline{T_{rd}} + T_{dt}) = 125.989s = 2 \text{ min } 5.989s$$

3. Šematski prikaz sistema dat je na sledećoj slici:



Ako intenzitet pristizanja zahteva u sistem obeležimo sa  $\lambda$ ,  $\lambda = \frac{1}{\bar{a}}$ ,  $\bar{a} = 16ms$ , a intenzitet obrade jednog kanala sa  $\mu$ ,  $\mu = \frac{1}{\bar{s}}$ ,  $\bar{s} = 25ms$ , tada dijagram stanja sistema izgleda kao na narednoj slici:



Stanje  $i$  predstavlja ono stanje sistema u kome u serveru postoji  $i$  zahteva.

Neka je  $\frac{\lambda}{\mu} = \rho$ . Balansne jednačine za ovaj sistem:

$$p_0 \cdot \lambda = p_1 \cdot \mu \Rightarrow p_1 = \frac{\lambda}{\mu} \cdot p_0 = \rho \cdot p_0$$

$$p_1 \cdot \lambda = p_2 \cdot 2\mu \Rightarrow p_2 = \frac{\lambda}{2\mu} \cdot p_1 = \frac{\rho^2}{2} \cdot p_0$$

$$p_2 \cdot \lambda = p_3 \cdot 2\mu \Rightarrow p_3 = \frac{\lambda}{2\mu} \cdot p_2 = \frac{\rho^3}{2^2} \cdot p_0$$

...

$$p_9 \cdot \lambda = p_{10} \cdot 2\mu \Rightarrow p_{10} = \frac{\lambda}{2\mu} \cdot p_9 = \frac{\rho^{10}}{2^9} \cdot p_0$$

$$\sum_i p_i = 1 \Rightarrow p_0 \cdot \left( 1 + \rho + \frac{\rho^2}{2} + \dots + \frac{\rho^{10}}{2^9} \right) = 1$$

$$p_0 \cdot \left( 1 + \rho \left( 1 + \frac{\rho}{2} + \left( \frac{\rho}{2} \right)^2 + \dots + \left( \frac{\rho}{2} \right)^9 \right) \right) = p_0 \cdot \left( 1 + \frac{\rho \left( 1 - \left( \frac{\rho}{2} \right)^{10} \right)}{1 - \frac{\rho}{2}} \right)$$

$$\rho = \frac{\lambda}{\mu} = \frac{\bar{s}}{\bar{a}} = 0.78125, \quad p_0 = 0.13266$$

$$p_1 = \rho \cdot p_0 = 0.2073$$

$$\text{Iskorišćenje servera: } U = 1 - \left( p_0 + \frac{1}{2} p_1 \right) = 0.7637$$

Srednji broj poslova u sistemu:

$$J = \sum_{i=0}^{\infty} i \cdot p_i = \sum_{i=1}^{10} i \cdot p_i = p_1 + \sum_{i=2}^{10} i \cdot p_i = p_0 \cdot \left( \rho + \sum_{i=2}^{10} i \cdot \frac{\rho^i}{2^{i-1}} \right) =$$

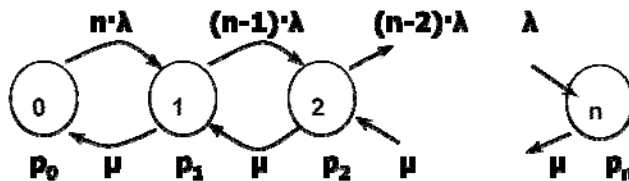
$$= p_0 \cdot \rho \cdot \sum_{i=1}^{10} i \cdot \frac{\rho^{i-1}}{2^{i-1}} = p_0 \cdot \rho \cdot \left( \frac{1 - \left(\frac{\rho}{2}\right)^{10}}{\left(1 - \frac{\rho}{2}\right)^2} - \frac{10 \left(\frac{\rho}{2}\right)^{10}}{1 - \frac{\rho}{2}} \right) = 3.1622$$

$$\text{Produktivnost: } X = \sum_{i=0}^{10} \lambda_i \cdot p_i = \lambda \cdot \sum_{i=0}^9 p_i = \lambda (1 - p_{10}) = \frac{1}{a} \left( 1 - p_0 \cdot \frac{\rho^{10}}{2^9} \right) = 61.095 \text{ posl / sec}$$

$$\text{Vreme odziva: } \bar{T} = \frac{J}{X} = 51.759 \text{ ms}$$

$$\text{Srednje vreme čekanja: } \bar{T}_q = \bar{T} - \bar{s} = 31.759 \text{ ms}$$

4. Dijagram stanja sistema prikazan je na slici:



Ako sa  $\lambda$  označimo intenzitet generisanja zahteva od strane jednog terminala ,

$\lambda = \frac{1}{\theta}$ , gde je  $\bar{\theta}$  srednje vreme razmišljanja korisnika (terminala) ,

a sa  $\mu$  označimo srednju brzinu procesora, odnosno intenzitet opsluživanja korisničkih zahteva, ( $\mu = \frac{1}{s}$ , gde je  $\bar{s}$  srednje vreme servisiranja zahteva ), tada su balansne jednačine za ovaj sistem:

$$p_0 \cdot n \cdot \lambda = p_1 \cdot \mu \Rightarrow p_1 = p_0 \cdot n \cdot \rho, \text{ gde je sa } \rho \text{ obeležen odnos } \frac{\lambda}{\mu} = \frac{1}{k}$$

$$p_1 \cdot (n-1) \cdot \lambda = p_2 \cdot \mu \Rightarrow p_2 = p_0 \cdot n \cdot (n-1) \cdot \rho^2$$

....

$$p_{n-1} \cdot \lambda = p_n \cdot \mu \Rightarrow p_n = p_0 \cdot n \cdot (n-1) \cdot \dots \cdot 1 \cdot \rho^n = p_0 \cdot n! \cdot \rho^n$$

Zbir svih verovatnoća je 1, pa dobijamo:

$$p_0 (1 + n\rho + n(n-1)\rho^2 + \dots + n! \rho^n) = 1$$

$$\frac{1}{p_0(n)} = 1 + n\rho + n(n-1)\rho^2 + \dots + n! \rho^n$$

Možemo uočiti da važi sledeća rekurzivna veza:

$$\frac{1}{p_0(n-1)} = 1 + (n-1)\rho + (n-1)(n-2)\rho^2 + \dots + (n-1)! \rho^{n-1}$$

$$\frac{1}{p_0(n)} = 1 + n\rho \cdot \frac{1}{p_0(n-1)}$$

Kako je iskorišćenje procesora jednako verovatnoći da je procesor zauzet, to je

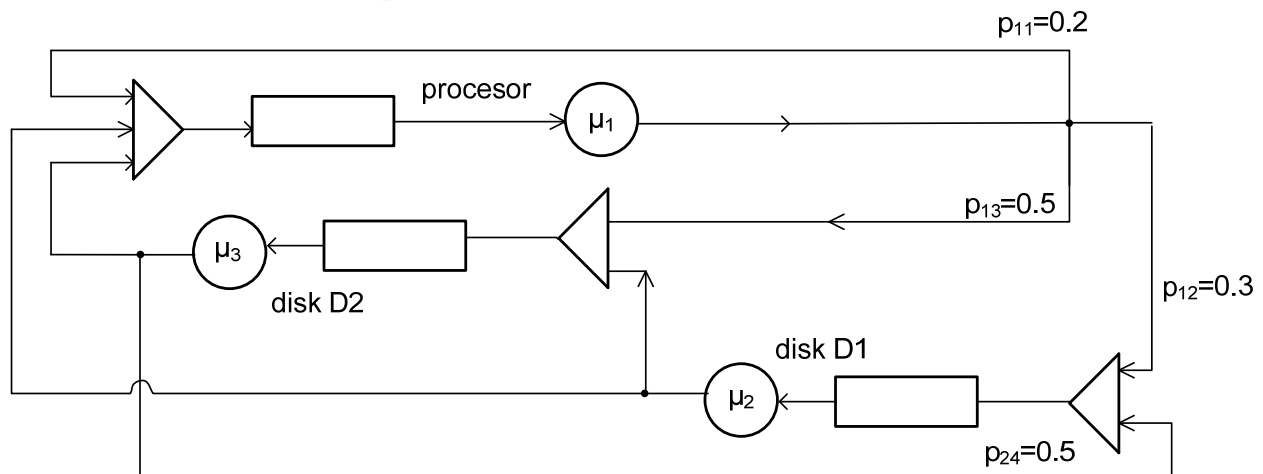
$$U(n)=1-p_0(n).$$

$$\frac{1}{p_0(n+1)} = 1 + (n+1)\rho \cdot \frac{1}{1-U(n)}$$

$$U(n+2) = 1 - p_0(n+2) = 1 - \frac{1}{1 + (n+2)\rho \cdot \frac{1}{p_0(n+1)}} = 1 - \frac{1}{1 + (n+2)\rho \cdot \left(1 + (n+1)\rho \cdot \frac{1}{1-U(n)}\right)}$$

$$U(n+2) = 1 - \frac{1}{1 + \frac{(n+2)}{k} \cdot \left(1 + \frac{(n+1)}{k(1-U)}\right)}$$

5. a) Dati sistem šematski je prikazan na narednoj slici.



$$\mu_1 = \frac{1}{s_p} = 100 \text{ sec}^{-1} \quad \mu_2 = \frac{1}{s_{D1}} = 25 \text{ sec}^{-1} \quad \mu_3 = \frac{1}{s_{D2}} = 40 \text{ sec}^{-1}$$

Verovatnoće  $p_{32}$  i  $p_{23}$  se mogu odrediti na sledeći način, znajući da je protok kroz oba diska isti (svaki zahtev koji prođe kroz jedan disk, proći će i kroz drugi):

Ako sa  $X$  obeležimo protok kroz procesor, tada je protok kroz svaki disk  $0.8X$ . Kako protok kroz granu 12 iznosi  $0.3X$ , tada će protok kroz granu 32 iznositi  $0.5X$ . Analogno će protok kroz granu 23 iznositi  $0.3X$ . Kako je protok kroz drugi disk  $0.8X$ , a protok kroz granu 23  $0.3X$ , tada je

$$p_{23} = \frac{0.3 \cdot X}{0.8 \cdot X} = \frac{3}{8}, \quad p_{21} = 1 - p_{23} = \frac{5}{8}. \text{ Analogno je:}$$

$$p_{32} = \frac{0.5 \cdot X}{0.8 \cdot X} = \frac{5}{8}, \quad p_{31} = 1 - p_{32} = \frac{3}{8}$$

Gordon-Newell-ove jednačine za data 4 resursa:

$$-(1-p_{11})\mu_1 x_1 + p_{21}\mu_2 x_2 + p_{31}\mu_3 x_3 = 0$$

$$p_{12}\mu_1 x_1 - (1-p_{22})\mu_2 x_2 + p_{32}\mu_3 x_3 = 0$$

$$p_{13}\mu_1 x_1 + p_{23}\mu_2 x_2 - (1-p_{33})\mu_3 x_3 = 0$$

Usvajajući da je  $x_1=1$ , dobijamo:

$$-0.8 \cdot 100 + \frac{5}{8} \cdot 25 \cdot x_2 + \frac{3}{8} \cdot 40 \cdot x_3 = 0$$

$$0.3 \cdot 100 - 25 \cdot x_2 + \frac{5}{8} \cdot 40 \cdot x_3 = 0$$

$$0.5 \cdot 100 + \frac{3}{8} \cdot 25 \cdot x_2 - 40 \cdot x_3 = 0$$

Rešenja su:  $x_1 = 1$ ,  $x_2 = 3.2$ ,  $x_3 = 2$

b)

#	X1=1	X2=3.2	X3=2
0	1	1	<b>1</b> = <b>G(0)</b>
1	1	4.2	<b>6.2</b> = <b>G(1)</b>
2	1	14.44	<b>26.84</b> = <b>G(2)</b>
3	1	47.208	<b>100.888</b> = <b>G(3)</b>
4	1	152.0656	<b>353.8416</b> = <b>G(4)</b>

$$g = \frac{G(3)}{G(4)} = 0.28512$$

Iskoriscenje procesora:  $U_p = g \cdot x_1 = g = 0.28512$

Iskoriscenje diska 1:  $U_{D1} = g \cdot x_2 = 0.91239$

Iskoriscenje diska 2:  $U_{D3} = g \cdot x_3 = 0.57024$

Protok kroz procesor:  $X_p = \frac{U_p}{s_p} = 28.51219 \text{ posl/sec}$

Protok kroz disk 1:  $X_{D1} = \frac{U_{D1}}{s_{D1}} = 22.80975 \text{ posl/sec}$

Protok kroz disk 2:  $X_{D2} = \frac{U_{D2}}{s_{D2}} = 22.80975 \text{ posl/sec}$

Vreme odziva:  $R = \frac{n}{X_p} = 140.29ms$

Usko grlo sistema je prvi disk, jer ima najveće iskorišćenje.

c) Srednji broj poslova u drugom disku:

$$\begin{aligned} \bar{n}_{D1} &= \sum_{j=1}^n x_3^j \cdot \frac{G(n-j)}{G(n)} = 2 \cdot \frac{G(3)}{G(4)} + 4 \cdot \frac{G(2)}{G(4)} + 8 \cdot \frac{G(1)}{G(4)} + 16 \cdot \frac{G(0)}{G(4)} = \\ &= \frac{16 \cdot G(0) + 8 \cdot G(1) + 4 \cdot G(2) + 2 \cdot G(3)}{G(4)} \approx 1.059 \end{aligned}$$